

Practical relevance of the state-of-the-art exact VRP solvers

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Plan of the talk

Exact SOTA algorithms for vehicle routing and their performance

POPMUSIC matheuristic for VRPs

VRPSolverEasy Python package

Contents

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Set-partitioning formulation

- ▶ Set of feasible routes = set P of resource constrained paths in graph $G = (V \cup \{0\}, A)$
- ▶ $h_a^p = 1$ if and only if path $p \in P$ contains arc a , otherwise 0
- ▶ Variable x_a — arc $a \in A$ is used in the solution or not
- ▶ Variable λ_p — path $p \in P$ is used in the solution or not

$$\text{Min } \sum_{a \in A} c_a x_a$$

$$\text{s.t. } \sum_{a \in \delta(v)} x_a = 2, \quad v \in V,$$

$$Bx \leq b, \quad (\text{add. constraints and robust cuts})$$

$$D\lambda \leq d, \quad (\text{non-robust cuts})$$

$$x_a = \sum_{p \in P} h_a^p \lambda_p, \quad a \in A,$$

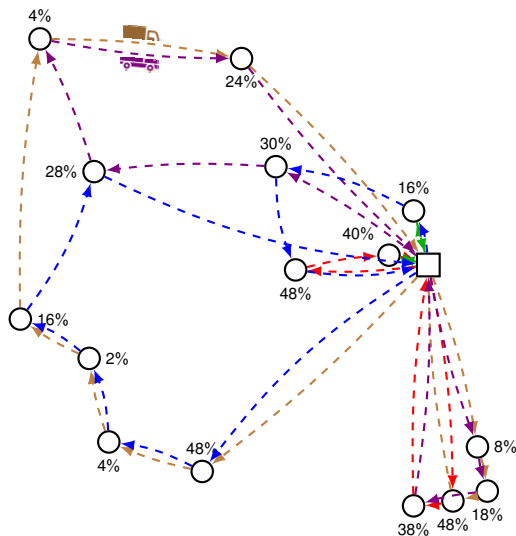
$$\sum_{p \in P} \lambda_p \leq K,$$

$$x_a, \lambda_p \in \{0, 1\}, \quad a \in A, p \in P.$$

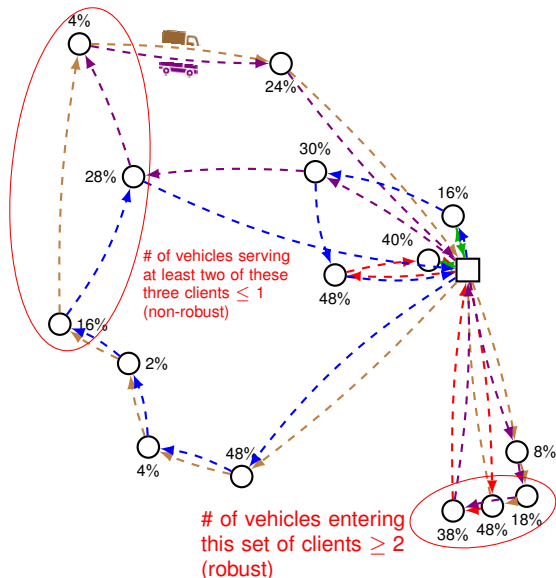
Column and cut generation: illustration

One **continuous variable** per feasible route.

Pricing problem is the **Resource Constrained Shortest Path** problem.



Column and cut generation: illustration



One continuous variable per feasible route.

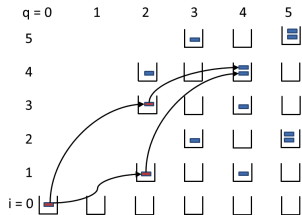
Pricing problem is the Resource Constrained Shortest Path problem.

Additional constraints (cuts) are added to reduce the integrality gap

Nodes (customers) are generalized to packing sets.

Resource constrained shortest path problem (RCSP)

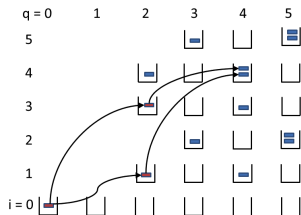
Labeling algorithm



- ▶ Enumeration of partial paths
- ▶ Relies heavily on domination
- ▶ Bi-directional search
- ▶ Buckets
- ▶ Completion bounds

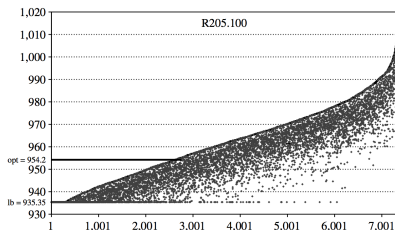
Resource constrained shortest path problem (RCSP)

Labeling algorithm



- ▶ Enumeration of partial paths
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Problem restriction using reduced cost arguments



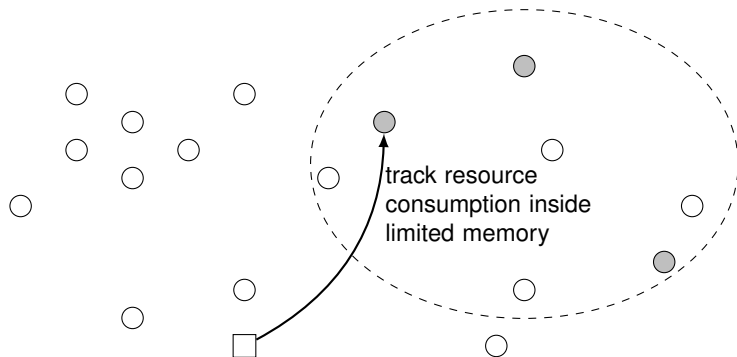
- ▶ Remove arcs contained only in paths with reduced cost $>$ primal-dual gap
- ▶ Enumerate all elementary paths with reduced cost \leq primal-dual gap

Figure from [Irnich et al., 2010]

Pricing problem relaxation: limited-memory concept

Specific RCSP instances

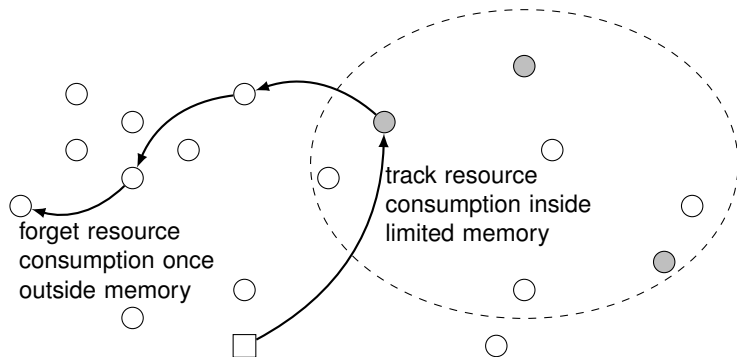
Few global resources but **hundreds of local resources**.



Pricing problem relaxation: limited-memory concept

Specific RCSP instances

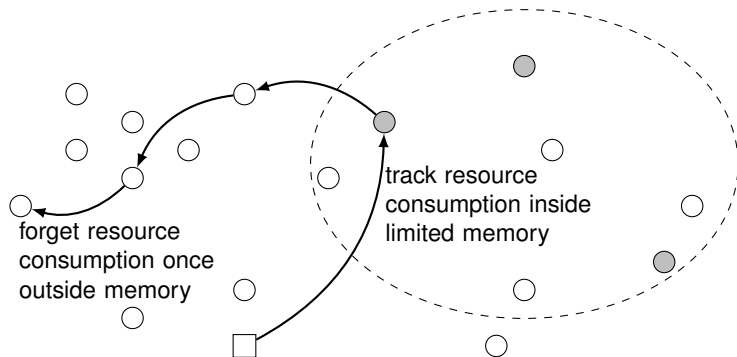
Few global resources but **hundreds of local resources**.



Pricing problem relaxation: limited-memory concept

Specific RCSP instances

Few global resources but **hundreds of local resources**.



Dynamic relaxation

Restrict relaxation after column generation convergence:
increase limited memories

Heuristics and branching

Generic heuristics: diving and restricted master IP

- ▶ **Work fine** only for instances with **short routes**
- ▶ **Cannot efficiently use diving heuristic** due to resource relaxation and a significant primal-dual gap.
- ▶ **Preliminary path enumeration with false gap** is required before running restricted master heuristic.

Branching

- ▶ **Strong branching** is important
- ▶ **Aggregated branching** may be important

Some history

- ▶ [Balinski and Quandt, 1964] set-partitioning formulation for CVRP
- ▶ [Laporte and Nobert, 1983] branch-and-cut, rounded capacity cuts
- ▶ [Desrosiers et al., 1984] first branch-and-price
- ▶ [Lysgaard et al., 2004] best branch-and-cut algorithm
- ▶ [Fukasawa et al., 2006] robust branch-cut-and-price
- ▶ [Baldacci et al., 2008] path enumeration technique
- ▶ [Jepsen et al., 2008] (non-robust) subset-row cuts
- ▶ [Baldacci et al., 2011] *ng*-route relaxation
- ▶ [Pecin et al., 2017] limited-memory technique, best branch-cut-and-price
- ▶ Latest survey: [Costa et al., 2019]
- ▶ [Pessoa et al., 2020] VRP generic model and exact solver

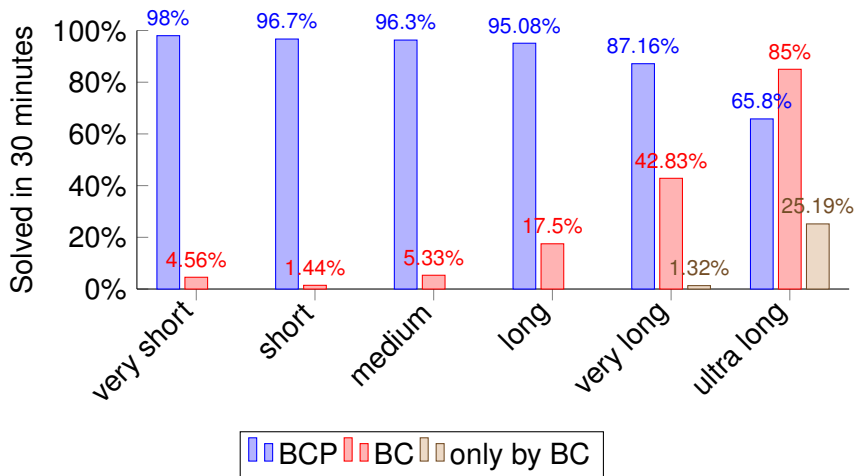
World record for the CVRP exact solving



Figure: Optimal solution for X-n865-k95 (solved in ≈ 20 days)

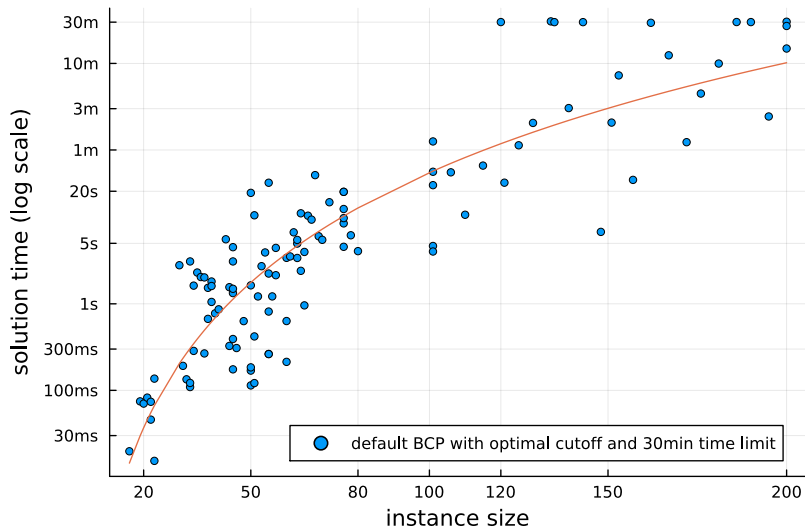
Computational comparison between BCP and BC

10'000 instances with 100 customers [Queiroga et al., 2022].



All 10'000 instances are solved to optimality thanks to cluster (aggregated) branching [Uchoa and Silva, 2022].

Solutions times (for CVRP)



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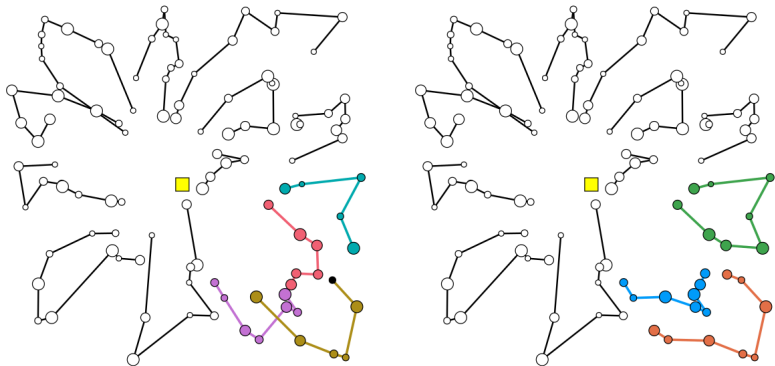
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POPMUSIC for CVRP: illustration

Partial **O**ptimization **M**etaheuristic **U**nder **S**pecial Intensification
Conditions [Taillard and Voss, 2002]



(a) Initial solution and a constructed subproblem. Seed client is marked in black.

(b) Improved solution after finding a better subsolution

POPMUSIC matheuristic for VRPs

Design choices

- ▶ Progressive increase of the subproblem size.
- ▶ Use of heuristic BCP
 - ▶ Time limit
 - ▶ False gap mechanism
 - ▶ Restricted master heuristic inside BCP

POPMUSIC matheuristic for VRPs

Design choices

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Using heuristic instead of BCP

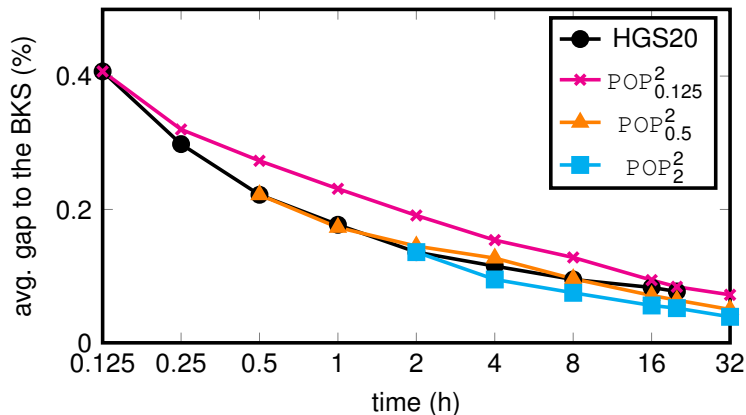
- ▶ We used [Vidal, 2022] heuristic instead of BCP for the DIMACS VRP challenge!
- ▶ Extensively studied in [Santini et al., 2023]
- ▶ However, **less generic** implementation.
- ▶ How to set the **termination criteria** for the subproblem?

Computational comparison for the CVRP

HGS20 : [Vidal, 2022].

POP_x^2 : POPMUSIC-BCP starting with x -hour solution of [Vidal, 2022].

X instances with $300 < n \leq 1000$ customers

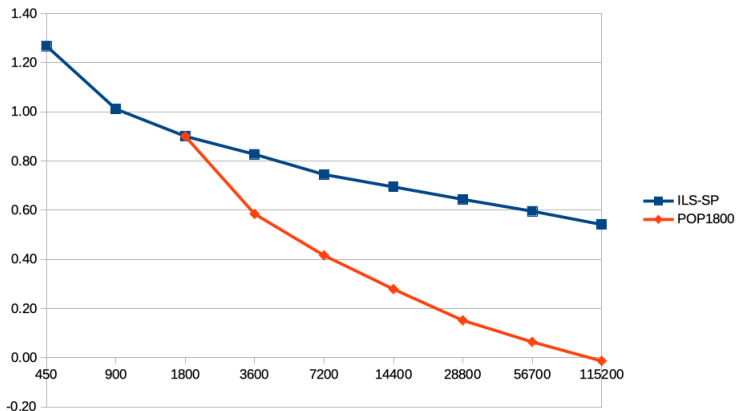


Computational results for the CVRP with backhauls

ILS-SP : [Subramanian and Queiroga, 2020].

POP1800: POPMUSIC-BCP starting with 30-minutes solution of [Subramanian and Queiroga, 2020]

Instances with 300–1000 clients.

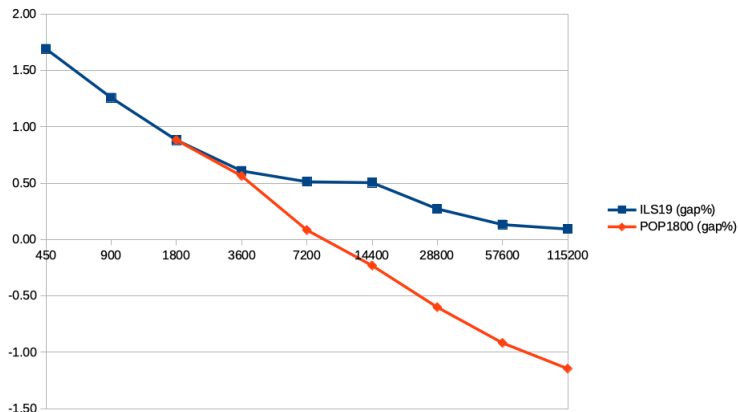


Computational results for the HFVRP

ILS19 : [Penna et al., 2019].

POP1800: POPMUSIC-BCP starting with 30-minutes solution of [Penna et al., 2019].

Instances of type XH with **300–1000 clients**, both limited and unlimited fleet.



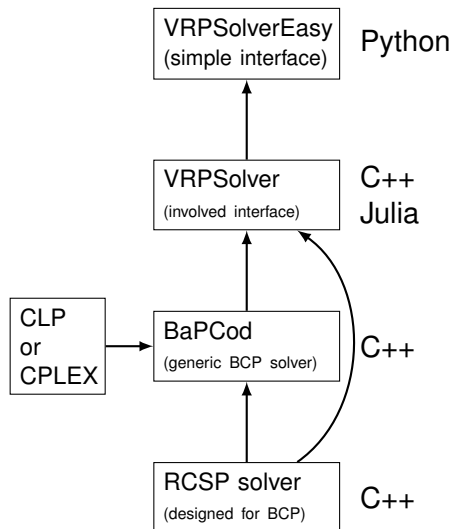
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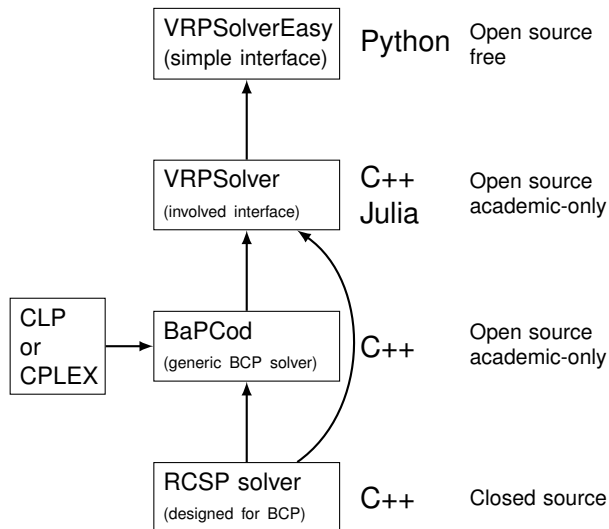
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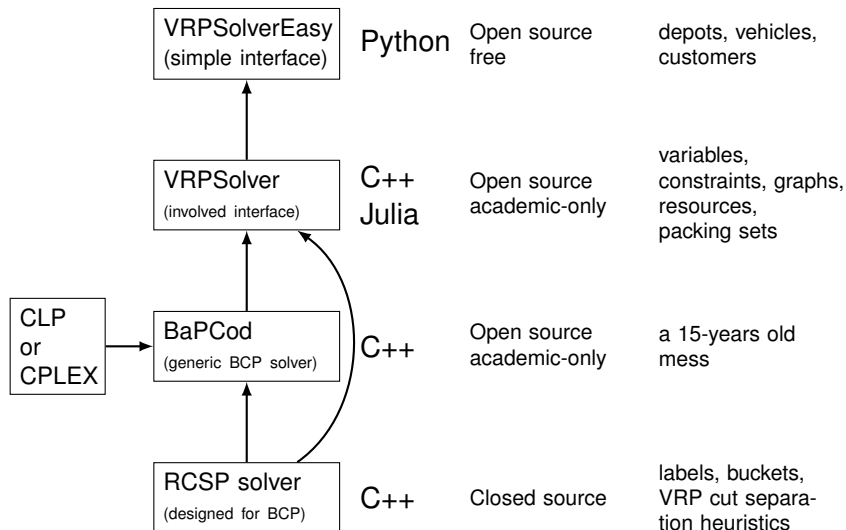
Four layers of VRPSolver implementation



Four layers of VRPSolver implementation



Four layers of VRPSolver implementation



VRPSolverEasy: compromise between generality and simplicity

<https://github.com/inria-UFF/VRPSolverEasy>

“OR-free” interface

Depots, customers, vehicles, links instead of variables, constraints, graphs, resources.

VRP variants covered

CVRP, VRPTW, HFVRP, MDVRP, OVRP, TOP, parallel links, site-dependent VRP, roaming delivery locations, and combinations.

VRP variants potentially covered

Arc routing, clustered VRP, generalized VRP, VRP with backhauls, Multi-trip VRP, Location-routing.

VRPSolverEasy: code example

```
1  from VRPSolverEasy.src import solver
2  model = solver.Model()
3  model.add_depot(id=0)
4  for i in range(data.nb_customers):
5      model.add_customer(id=i+1, demand=data.cust_demands[i])
6  for i,cust_i in enumerate(data.cust_coordinates):
7      for j in range(i + 1, len(data.cust_coordinates)):
8          dist = euclidean_distance(cust_i[0], cust_i[1],
9                                    data.cust_coordinates[j][0],
10                                   data.cust_coordinates[j][1])
11         model.add_link(start_point_id=i+1,
12                        end_point_id=j+1,
13                        distance=dist)
14     model.add_vehicle_type(id=1, start_point_id=0,
15                            end_point_id=0,
16                            max_number=data.nb_customers,
17                            capacity=data.vehicle_capacity,
18                            var_cost_dist=1)
19     model.solve()
20     if model.solution.is_defined():
21         print(model.solution)
```

VRPSolverEasy: computational results

Initial solution by OR-Tools run for $n/2$ sec. (included in solution time)

VRPSolverEasy time limit is 30 minutes

CVRP-Small, VRPTW-Solomon, HFVRP-Classic: ≤ 100 customers

gap_I, **gap_F** — OR-Tools gap, VRPSolverEasy gap

Problem	Dataset	Without built-in heuristic				With built-in heuristic			
		#opt	Time	gap _I	gap _F	#opt	Time	gap _I	gap _F
CVRP	Small	87/87	33.8	0.99	0.00	87/87	31.0	0.99	0.00
	X _≤ 200	14/26	844.9	4.36	2.35	16/26	617.3	4.36	1.96
VRPTW	Solomon	53/56	109.0	14.20	0.14	54/56	112.2	14.20	0.10
	HG-200	27/60	881.0	9.02	4.63	28/60	879.7	9.16	4.78
HFVRP	Classic	31/40	234.5	10.35	0.95	36/40	154.2	10.32	0.34
	XH _≤ 200	2/22	1730.4	18.07	13.44	8/22	1088.6	18.65	8.25
Total		215/291				229/291			

Practical relevance of exact VRP solvers

Possible now

- ▶ Finally easy to use!
- ▶ Small and moderate size (up to ≈ 100 customers)
- ▶ Not so long routes (up to ≈ 15 customers per route)
- ▶ For these instances
 - ▶ More generic than SOTA heuristic VRP solvers
 - ▶ More efficient than generic VRP solvers (OR-Tools, LocalSolver)

Perspectives

- ▶ **Industrial-quality** codes are needed!
- ▶ **Parallel mathheuristics** for larger instances
- ▶ Hybrid heuristics which **use dual information** from column and cut generation

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